Parallellizing large search in BOINC: a case study

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Brute-force search

Sometimes we just don’t have good algorithms.

- Satisfaction problem and constraint programming
- Combinatorial optimization
- Artificial intelligence
- Conjecture verification
- etc...

Quotation

When in doubt, use brute force.

– Ken Thompson
Backtracking
A **weird** number is a notion related to proper divisors. It is unknown whether an odd weird number exists, and the great mathematician Paul Erdös offered a prize for it.

\[ S(n) = \{ k \mid k < n, k \mid n \}. \sum_{k \in S(n)} k > n \land \forall S' \subset S(n), \sum_{k \in S'} k \neq n. \]

We want to search for an odd weird number.
Our case: Odd Weird Search

To check if a number is weird, we need to obtain its proper divisors.

- **Naïve way**
  For each odd number, we factorize it, calculate its proper divisors, solve a subset sum problem to determine if it is weird. Factorization very expensive, but easy to partition.

- **Backtracking**
  We construct (or search) directly the factorization, and backtrack whenever possible. The rest is the same. Much faster, but difficult to partition.

Of course we go the difficult but faster way.
Difficulties

Although much faster, backtracking is difficult to parallelize.

- No good and simple estimate of search volume, only rough and simple ones
- Irregularity of subtrees

In the BOINC context, it is even harder. We want to meet the following demand to please our volunteers:

- Reasonable and consistent runtime for every workunit (< 12h)
- Reasonable progress bar
- Easy workunit generation

None of them is easy.
Previous attempts

There are already some projects that parallelize their search of similar flavor.

- Rectilinear Crossing Number (with occasional extremely huge workunits)
- NQueens@home (smooth search space)
- SAT@home (expensive Monte Carlo estimation)
- etc...
No good and simple estimate, so we just use rough and simple ones! Rough estimate + Irregular subtrees ⇒ workunits of varying size

Our trick

Just force every workunit to stop after some time and then send back its checkpoint. We then recycle checkpoints sent back as workunits.

We wrote a mechanism to roughly estimate running time by counting operations in expensive functions. This estimate can be off by 20%, but still consistent, and we have a control. And we have a progress bar for free!
Our scheme

What does it look like?

For simplicity, three types of file share the same format.
Our scheme

Issues

At the end, throughput will drop due to low “liquidity”.

- Shorten deadline to increase “liquidity”
- Send more initial replica to shorten waiting time
- Compute locally when not many are left

To assure correctness, we do a quorum 2.
Our scheme only works when we only care about the search as a whole, and when the checkpoint is not large.
Conclusion

Our scheme can be used to parallelize a large class of search with a rough search volume estimate. In fact, an upperbound would work. In fact, it works as a potential solution to large workunits, if latency is not important.

Future work:

- Better recycle strategy
  - Estimation of search volume of recycled workunits
  - Priority of different recycled workunits

- Automatic control of “liquidity”

- Quantitative analysis?
Thanks for your attention!